

# Evaluation and Improvement of Generic-Emulating DPA Attacks

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# Outline

- 1. Background: generic-emulating DPA**
- 2. Two new generic-emulating distinguishers**
- 3. Improvement using cross-validation**
- 4. Experimental results**

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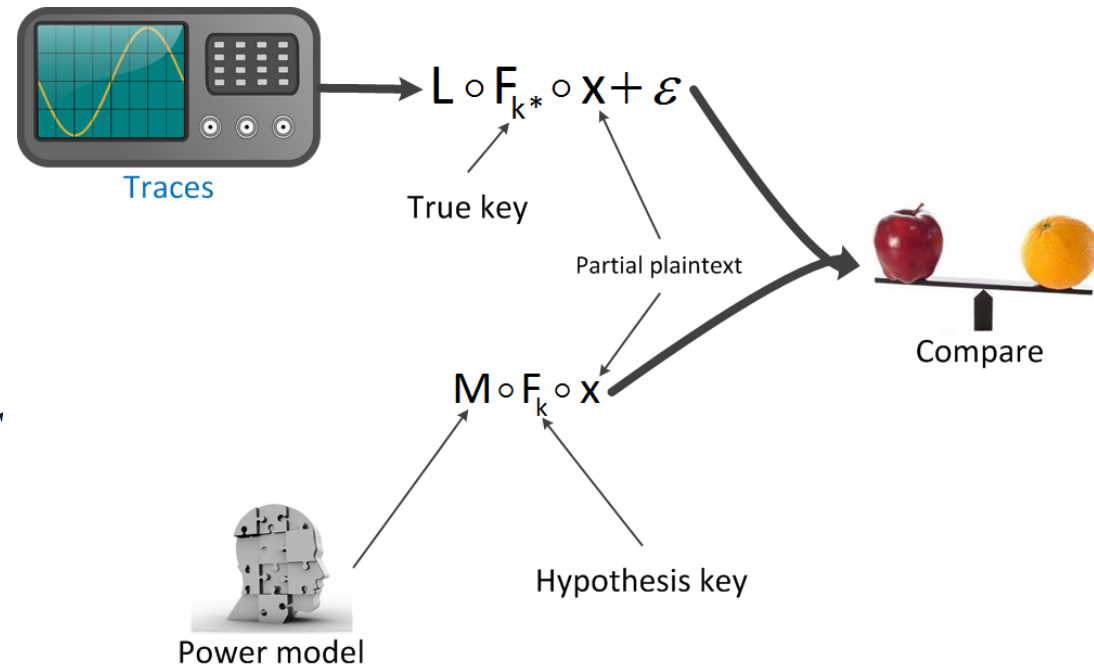
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- ◆ Schindler W, Lemke K, Paar C. A stochastic model for differential side channel cryptanalysis. CHES 2005.
- ◆ Doget J, Prouff E, Rivain M, et al. Univariate side channel attacks and leakage modeling. Journal of Cryptographic Engineering, 2011.
- ◆ Whitnall C, Oswald E, Standaert F X. The Myth of Generic DPA... and the Magic of Learning. CT-RSA 2014.

# 1.1 Differential Power Analysis (DPA)

## ❖ Problems of DPA:

- Choice of power model depends on the experiences of attacker
- *The impact of power variability is becoming more and more significant, which makes common power models much less respected in practice.*



## ❖ Solution:

- Generic DPA (e.g. MIA)

## 1.2 Generic DPA

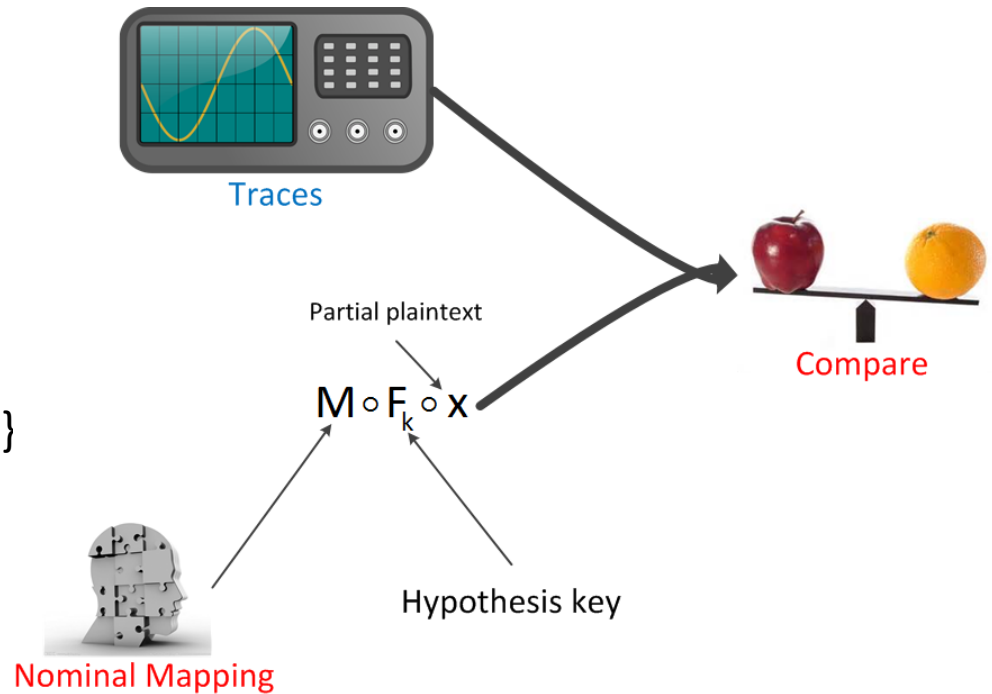
❖ Generic DPA use the nominal mapping as power model.

- We call the function  $M(\cdot)$  as nominal mapping if we have:

$$\{z \mid M(z) = M(z')\} \approx \{z \mid L(z) = L(z')\}$$

❖ Limitation of generic DPA:

- It doesn't work when the target function  $F_k(x)$  is injective (AES sbox)



## 1.3 The Power Model using Algebra Normal Form

❖ Fact: any real valued leakage function can be represented in algebra normal form (ANF).

❖ For Example:

Let  $z = (z_1, z_2, z_3)$  in  $\text{GF}(2)^3$

For any leakage function  $L(\cdot)$ , we have:

$$L(z) = \alpha_0 + \boxed{\alpha_1 z_1 + \alpha_2 z_2 + \alpha_3 z_3} + \boxed{\alpha_4 z_1 z_2 + \alpha_5 z_1 z_3 + \alpha_6 z_2 z_3} + \boxed{\alpha_7 z_1 z_2 z_3},$$

where  $\alpha_0$  to  $\alpha_7$  are the real number coefficients

terms of degree 1

terms of degree 2

terms of degree 3

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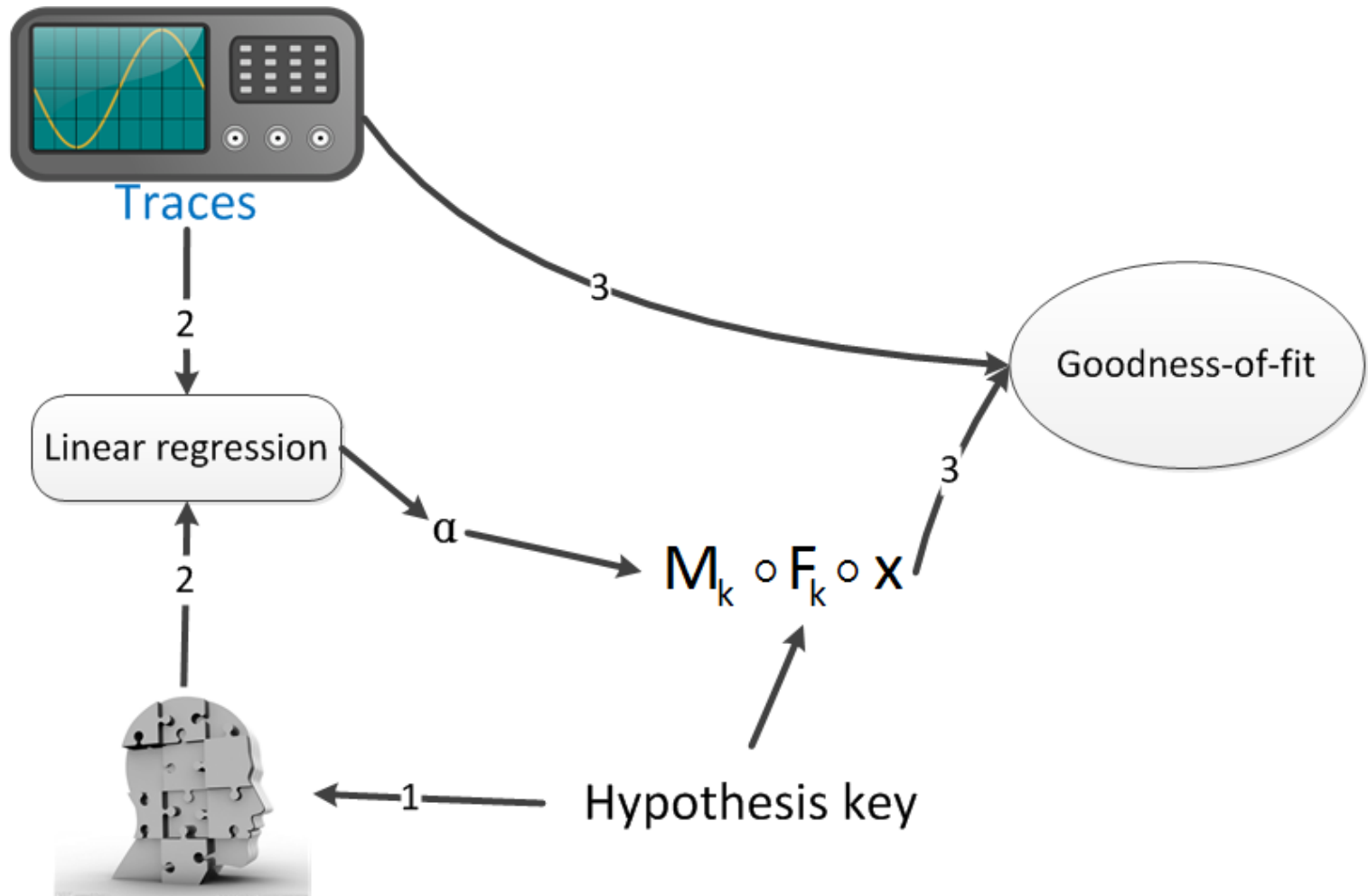
terms of degree 1

terms of degree 2

terms of degree 3

❖ Therefore, we can construct the nominal mapping power model using ANF

# 1.4 Linear Regression(LR)-based DPA

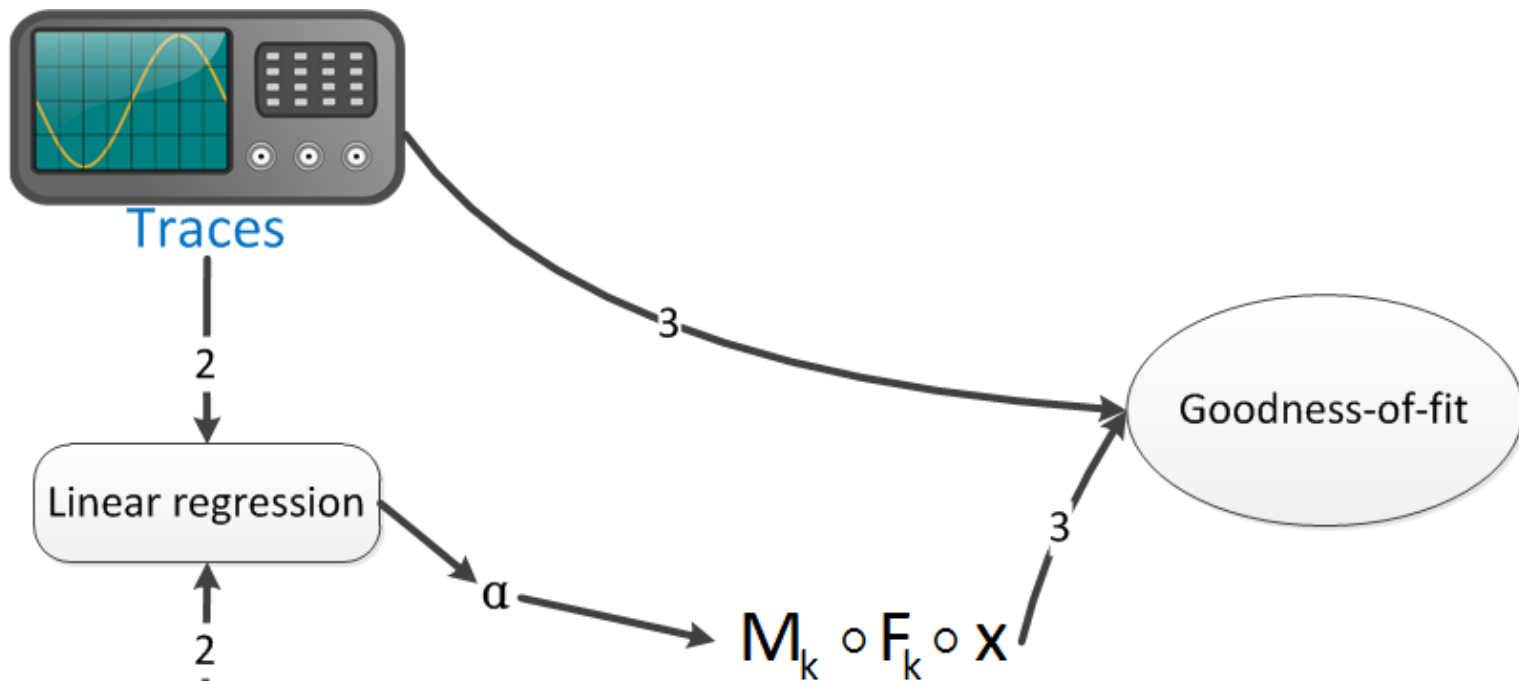


Power model: ANF

$$M_k(Z_{i,k}) = \alpha_0 + \sum_{u \in U} \alpha_u Z_{i,k}^u$$



# 1.4 Linear Regression(LR)-based DPA



**It is generic DPA**

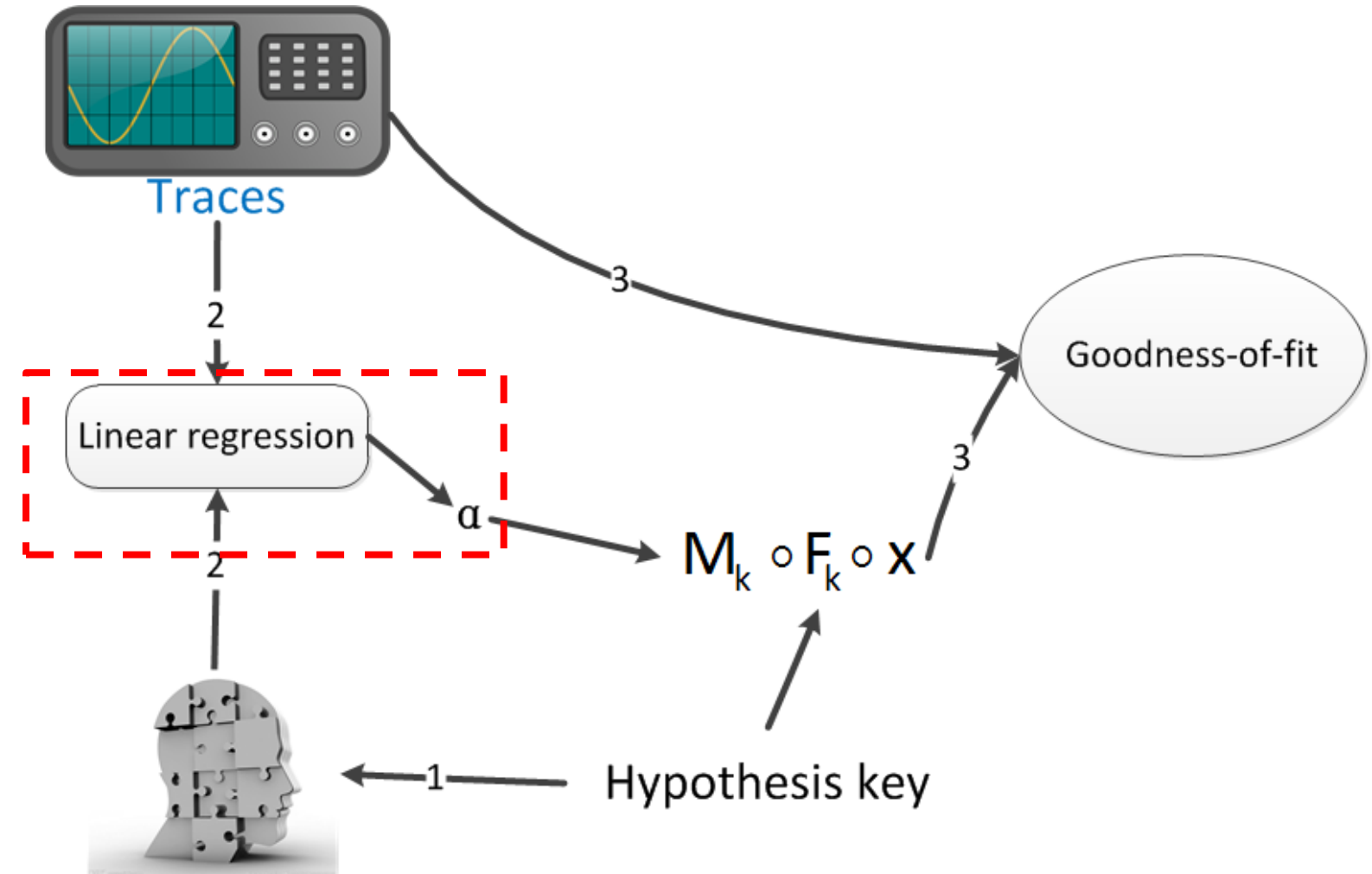
**And it doesn't work**

**with injective target function**

Power model: ANF

$$M_k(Z_{i,k}) = \alpha_0 + \sum_{u \in U} \alpha_u Z_{i,k}^u$$

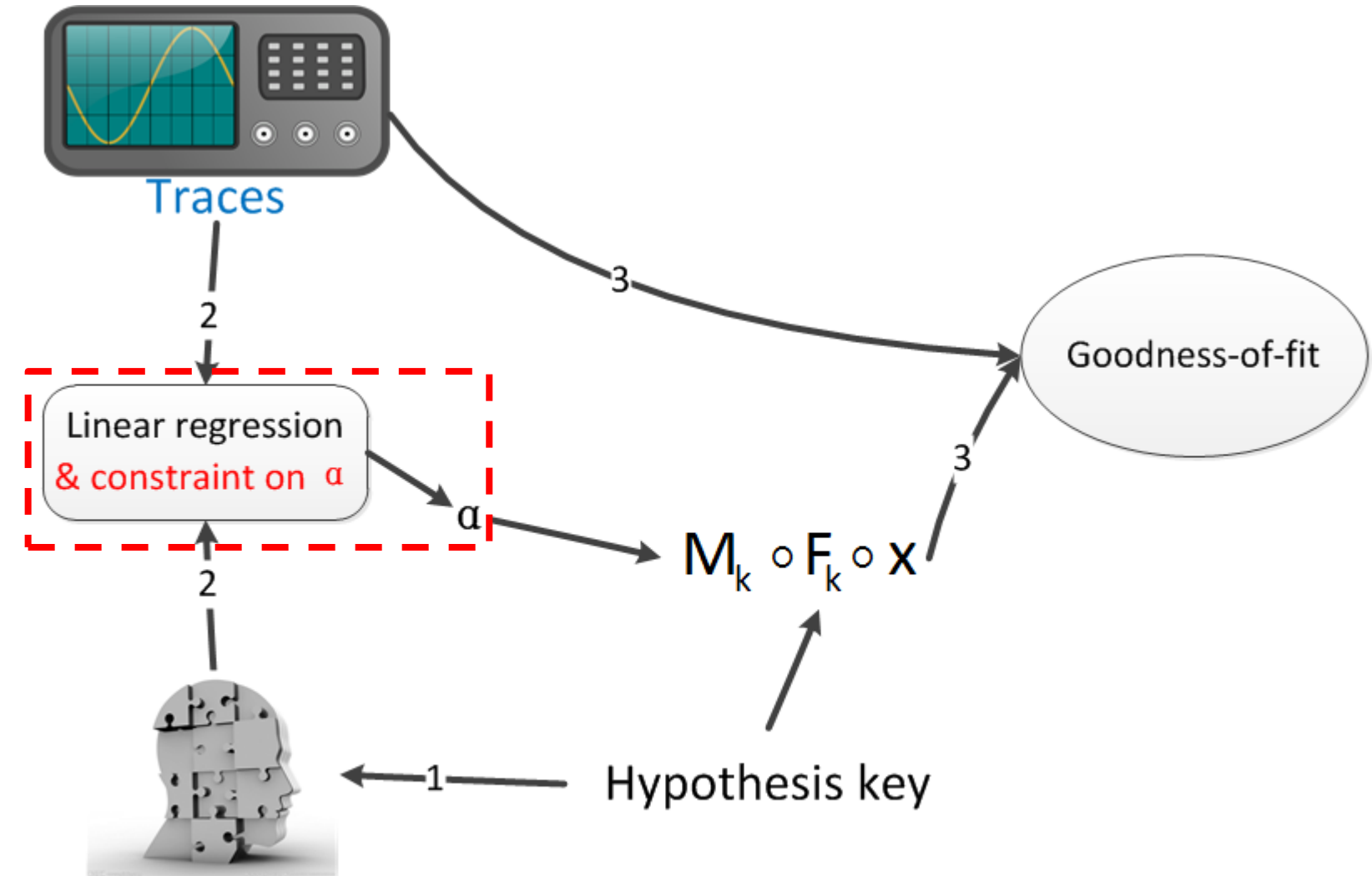
# 1.5 Generic-emulating DPA



Power model: ANF

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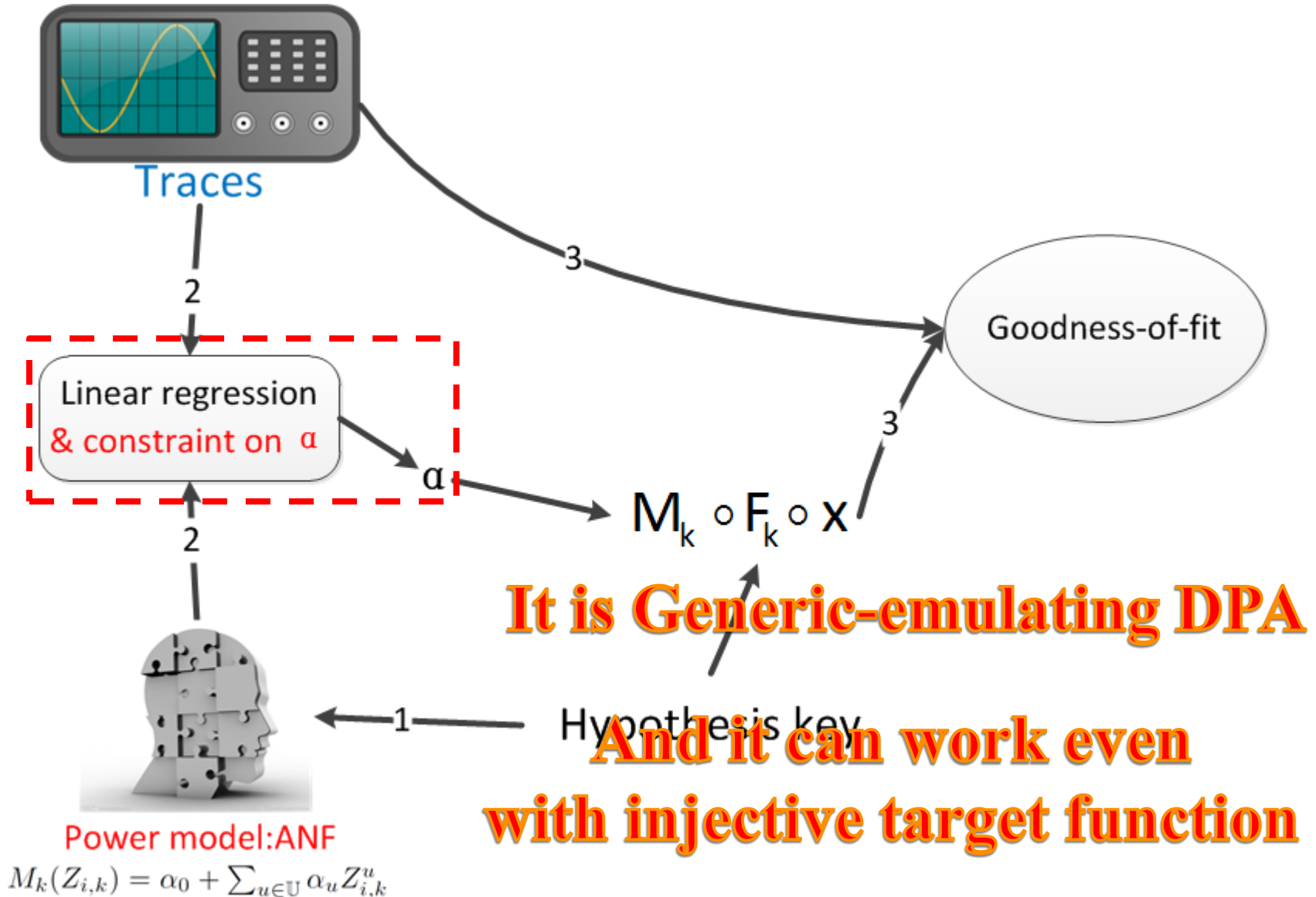
# 1.5 Generic-emulating DPA



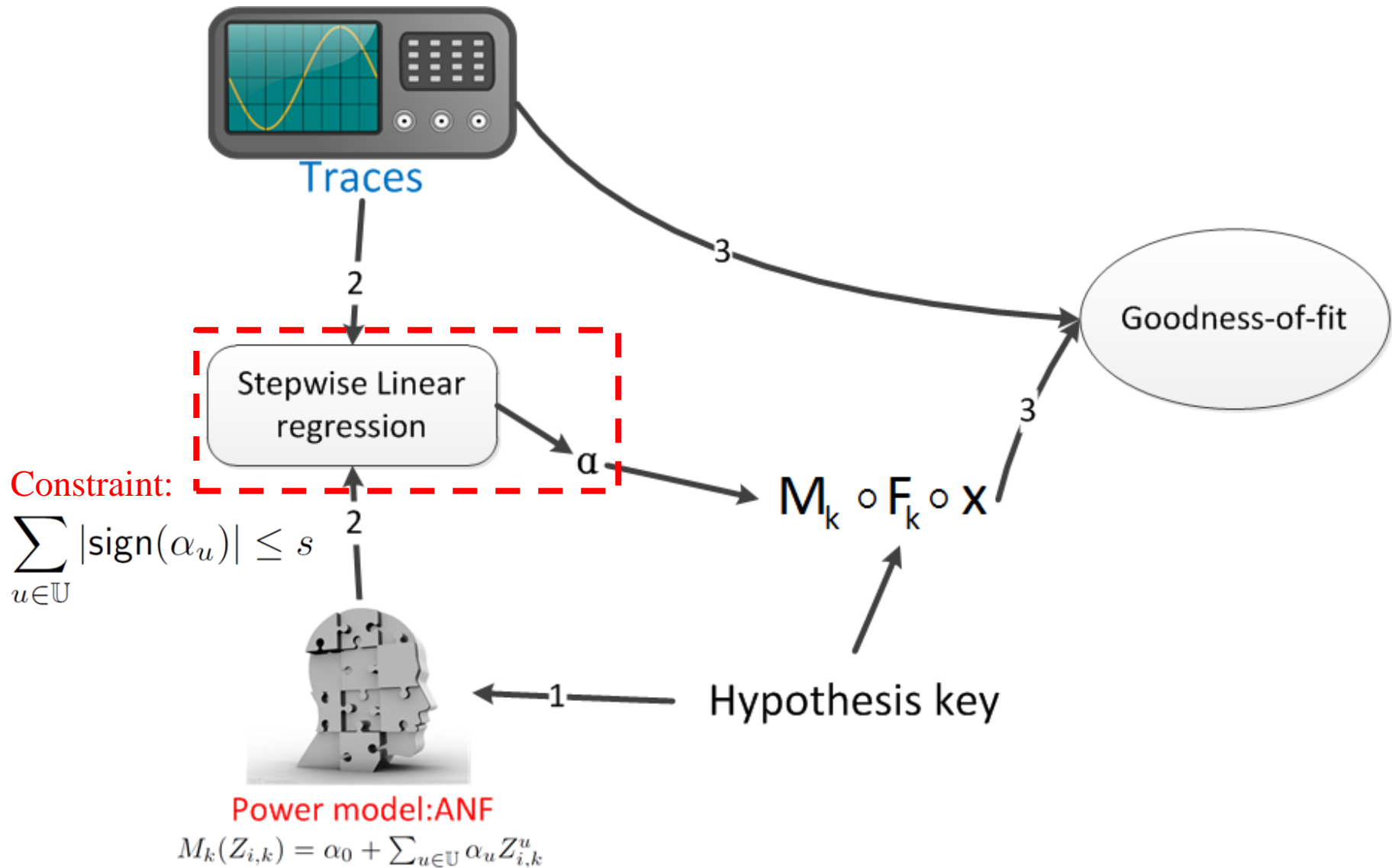
Power model: ANF

$$M_k(Z_{i,k}) = \alpha_0 + \sum_{u \in U} \alpha_u Z_{i,k}^u$$

# 1.5 Generic-emulating DPA



# 1.6 Stepwise Linear Regression (SLR)-based DPA



## 1.6 SLR-based DPA

- ❖ The coefficients in the leakage function are sparse
- ❖ Formal description:

$$\hat{\alpha}^{SLR} \stackrel{\text{def}}{=} \underset{\alpha}{\operatorname{argmin}} \sum_{i=1}^N (T_i - M_k(Z_{i,k}))^2$$

subject to  $\sum_{u \in \mathbb{U}} |\operatorname{sign}(\alpha_u)| \leq s$

# Outline

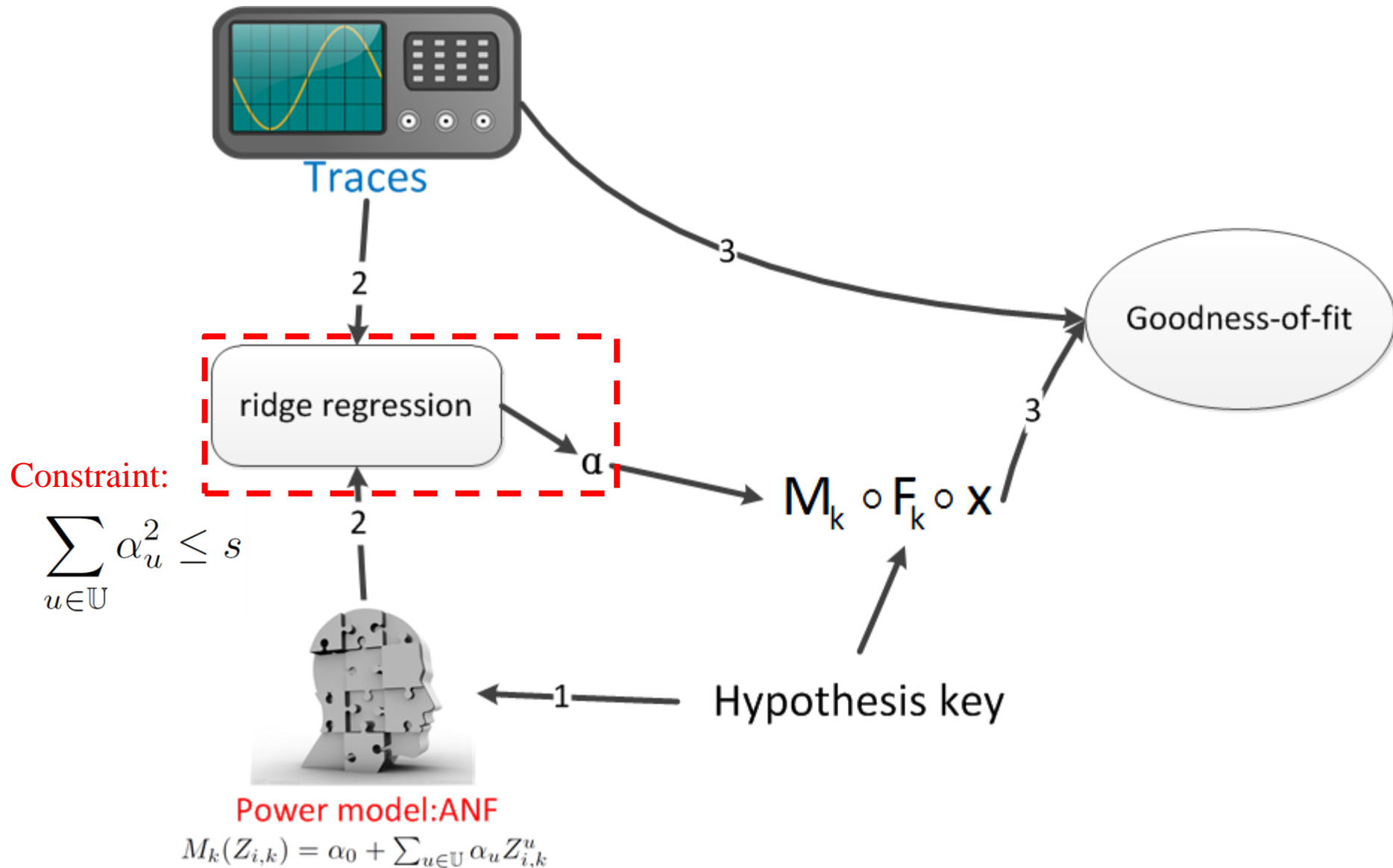
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2. **Two new generic-emulating distinguishers**
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## Motivation

- ❖ Two drawbacks in SLR-based DPA
  - Unstable outcomes in the high-noise regime
    - the insignificant coefficients are discarded, which makes the unstable outcomes
  - Less-satisfactory performance especially on real smart cards



## 2.1 Ridge-based Distinguishers



## 2.1 Ridge-based Distinguishers

- ❖ Ridge-based distinguisher shrinks coefficients by explicitly imposing an overall constraint on their size:

$$\hat{\alpha}^{ridge} \stackrel{\text{def}}{=} \underset{\alpha}{\operatorname{argmin}} \sum_{i=1}^N \left( T_i - M_k(Z_{i,k}) \right)^2$$

subject to  $\sum_{u \in \mathbb{U}} \alpha_u^2 \leq s$

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subject to  $\sum_{u \in U} \alpha_u^2 \leq s$

- ❖ An equivalent formulation:

$$\hat{\alpha}^{ridge} = \underset{\alpha}{\operatorname{argmin}} \left( \sum_{i=1}^N (T_i - M_k(Z_{i,k}))^2 + \lambda \sum_{u \in U} \alpha_u^2 \right)$$

## 2.1 Ridge-based Distinguishers

❖ The optimal solution is given by:

$$\hat{\alpha}^{ridge} = (\mathbf{U}_k^\top \mathbf{U}_k + \lambda \mathbf{I})^{-1} \mathbf{U}_k^\top T$$

where  $\mathbf{U}_k = (Z_{i,k}^u)_{i \in \{1,2,\dots,N\}, u \in \mathbb{F}_2^m \setminus \{0\}}$

$$\hat{\alpha}^{LR} = (\mathbf{U}_k^\top \mathbf{U}_k + \quad )^{-1} \mathbf{U}_k^\top T$$

shrink

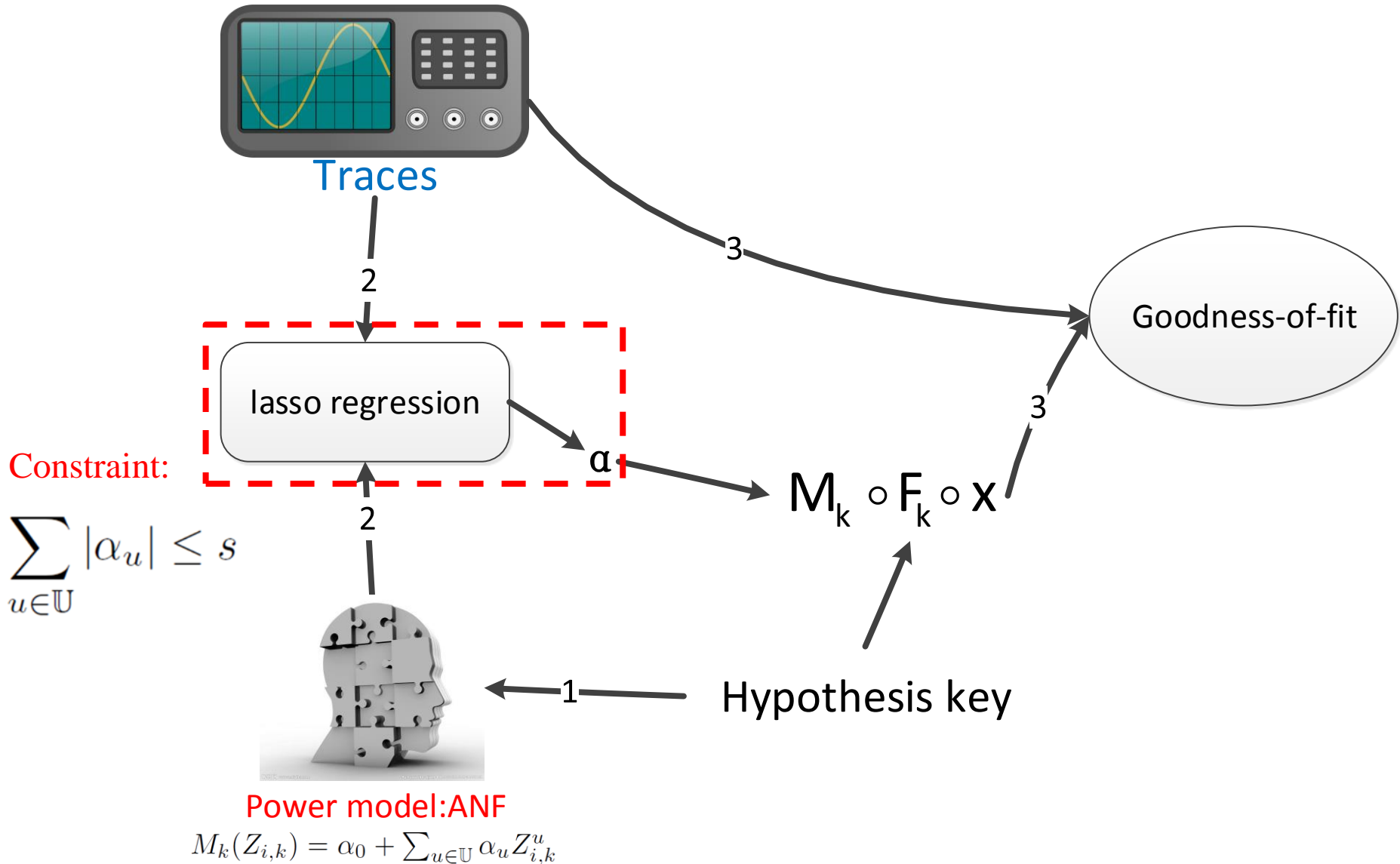
$$\hat{\alpha}^{ridge} = (\mathbf{U}_k^\top \mathbf{U}_k + \lambda \mathbf{I})^{-1} \mathbf{U}_k^\top T$$

## 2.2 How The Coefficients Shrink in Ridge-based Distinguishers

Amount of shrinkage of terms' coefficients ← Proportional to → Degrees of the terms

Consistent with leakage functions in practice

## 2.3 Lasso-based Distinguishers



## 2.3 Lasso-based Distinguishers

- ❖ The lasso-based distinguisher is similar to the ridge-based one excepted for a different constraint:

$$\hat{\alpha}^{lasso} \stackrel{\text{def}}{=} \underset{\alpha}{\operatorname{argmin}} \sum_{i=1}^N \left( T_i - M_k(Z_{i,k}) \right)^2$$

subject to  $\sum_{u \in \mathbb{U}} |\alpha_u| \leq s$

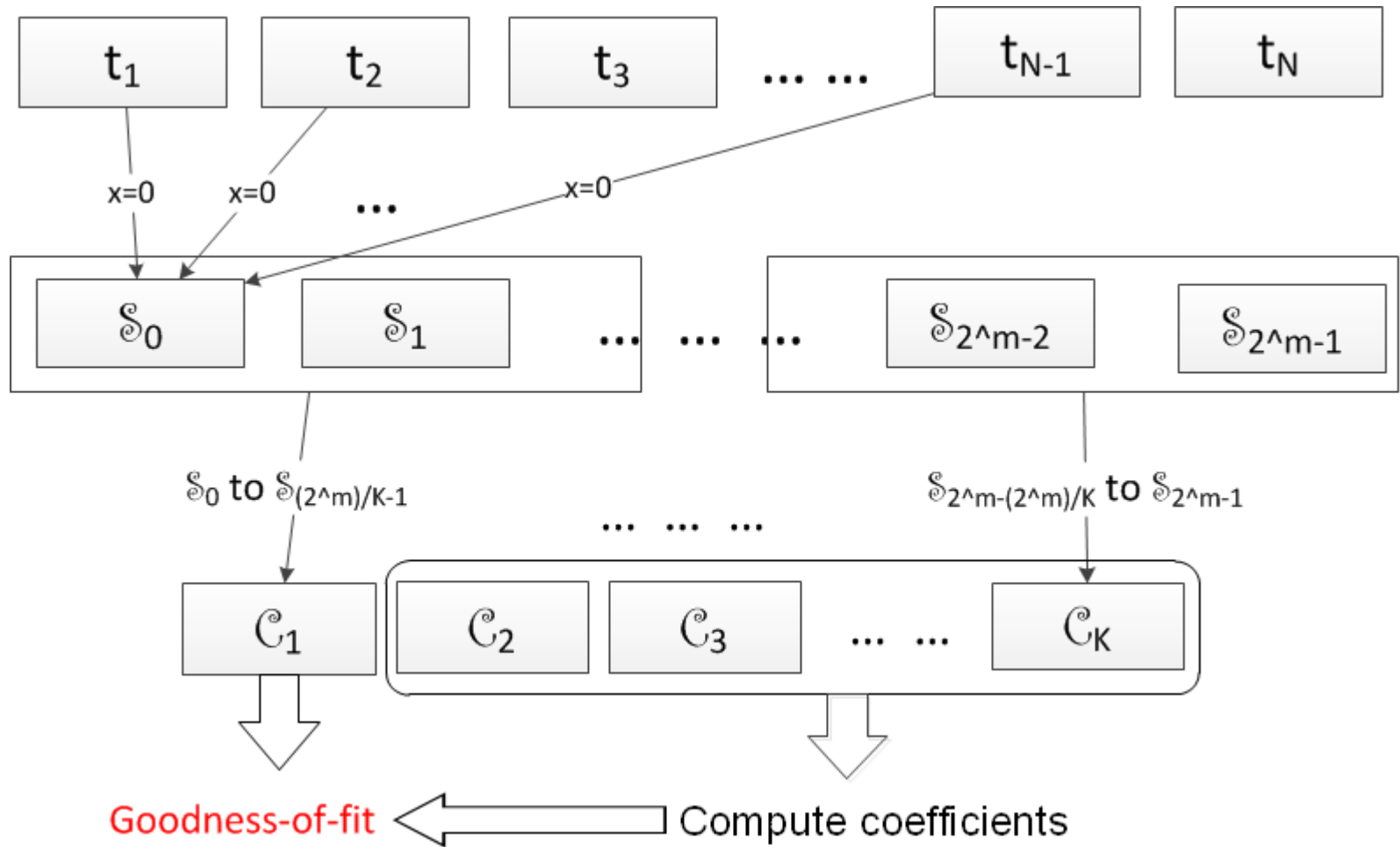
- ❖ Finding the optimal solution for lasso-based distinguishers is essentially a quadratic programming problem

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3. **Improvement using cross-validation**
4. Experimental results



# 3 Cross-validation

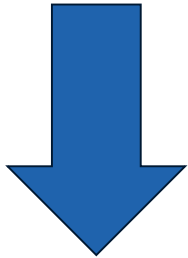


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# 4.1.1 SLR-based Distinguisher is Not Stable

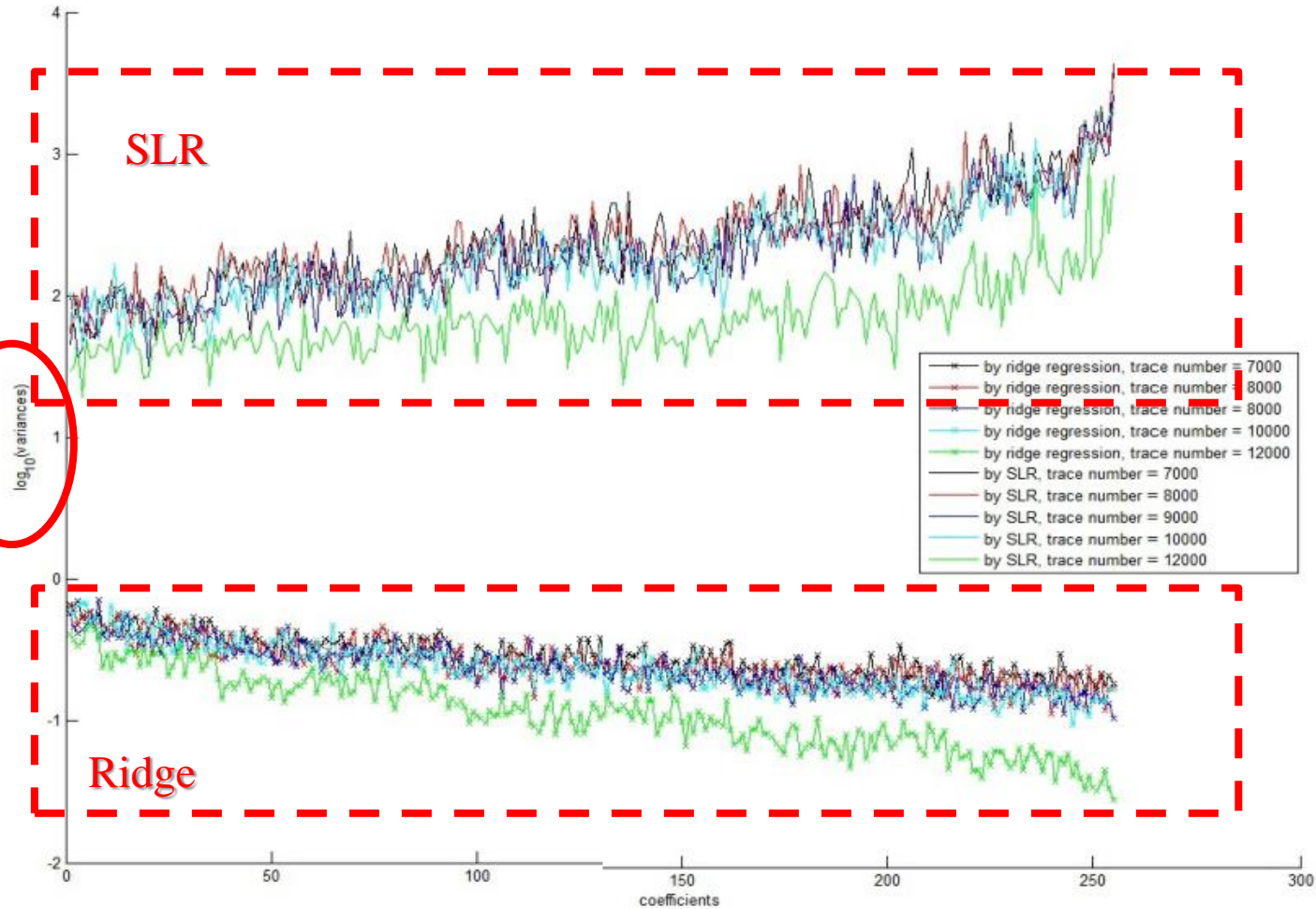
High variance  
of coefficients



Unstable

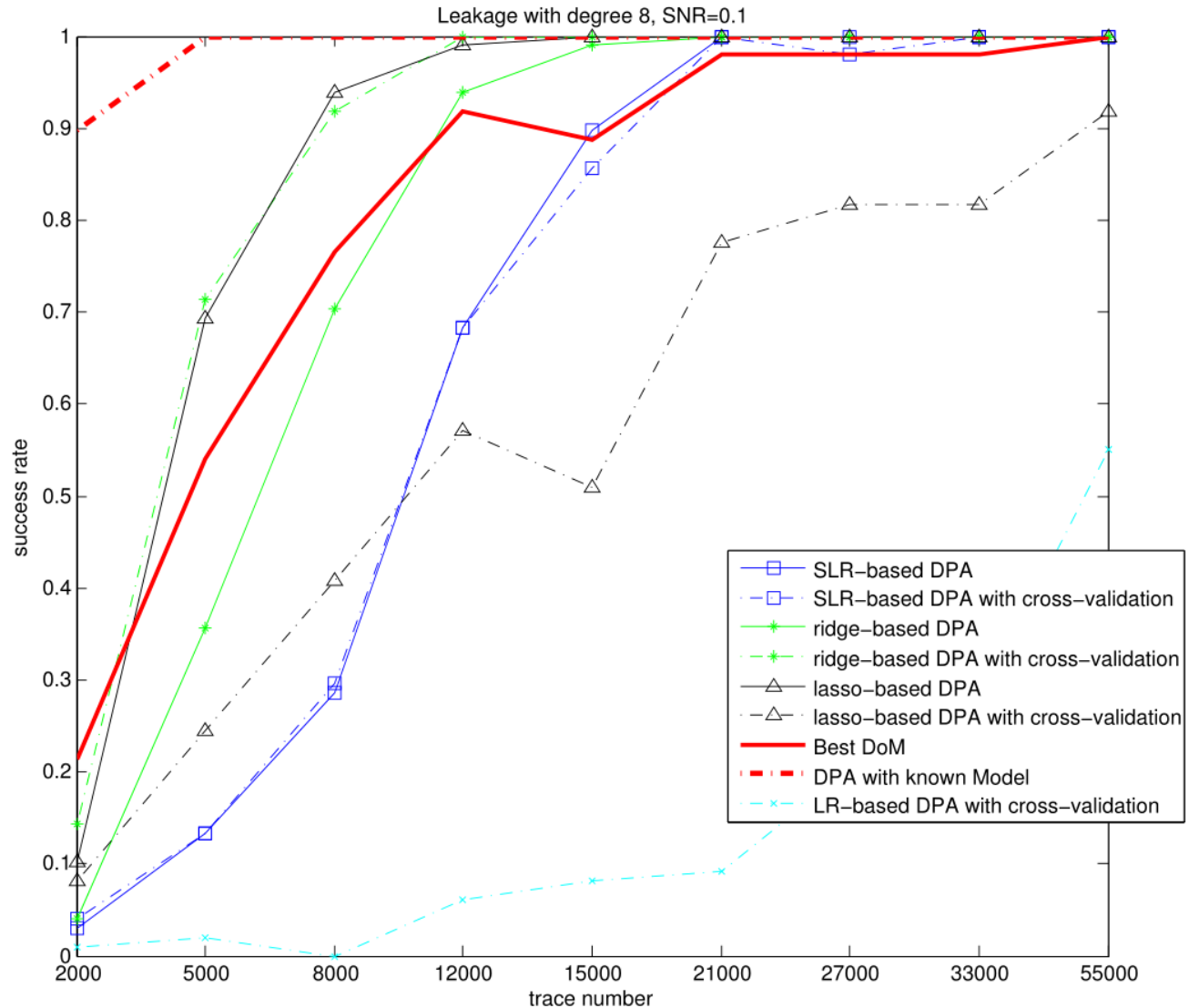


Logarithmic



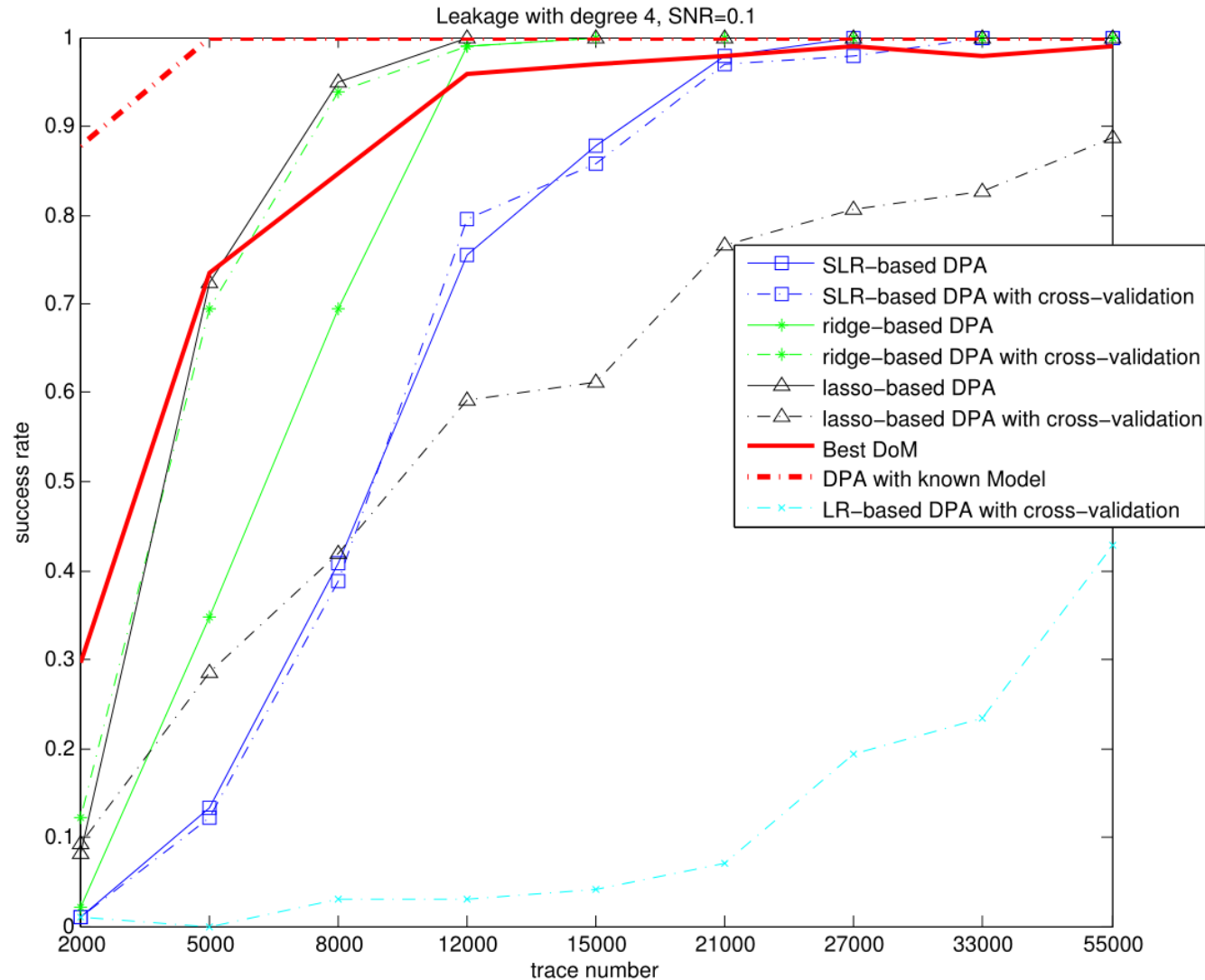
## 4.1.2 A Comparison of Various Attacks

- Leakage with degree 8
- Ridge-based DPA with C-V and lasso-based DPA are best
- New generic-emulating DPAs perform better than SLR-based One
- C-V improves the ridge-based DPA



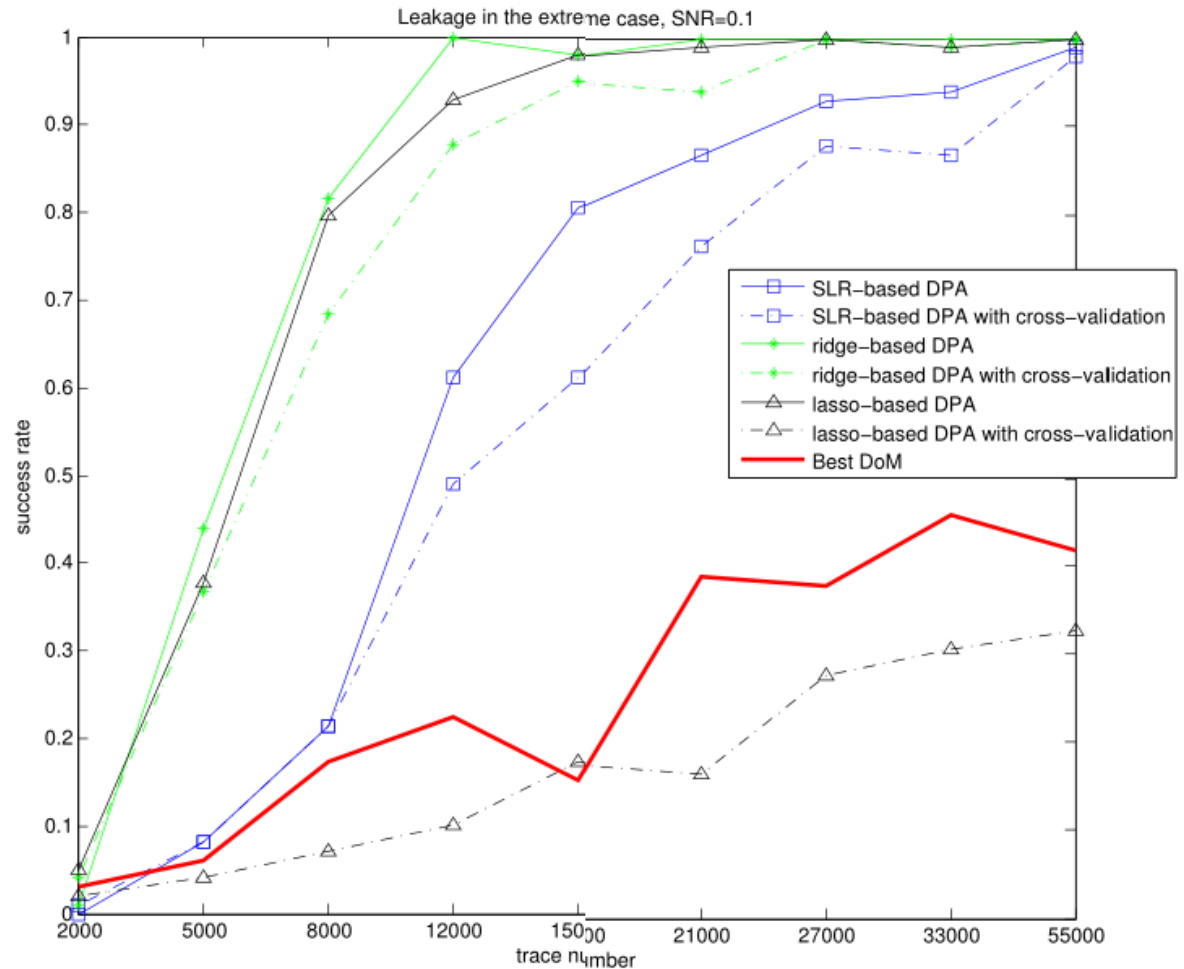
## 4.1.2 A Comparison of Various Attacks

- Leakage with degree 4
- The Best DoM becomes better in lower degree leakage



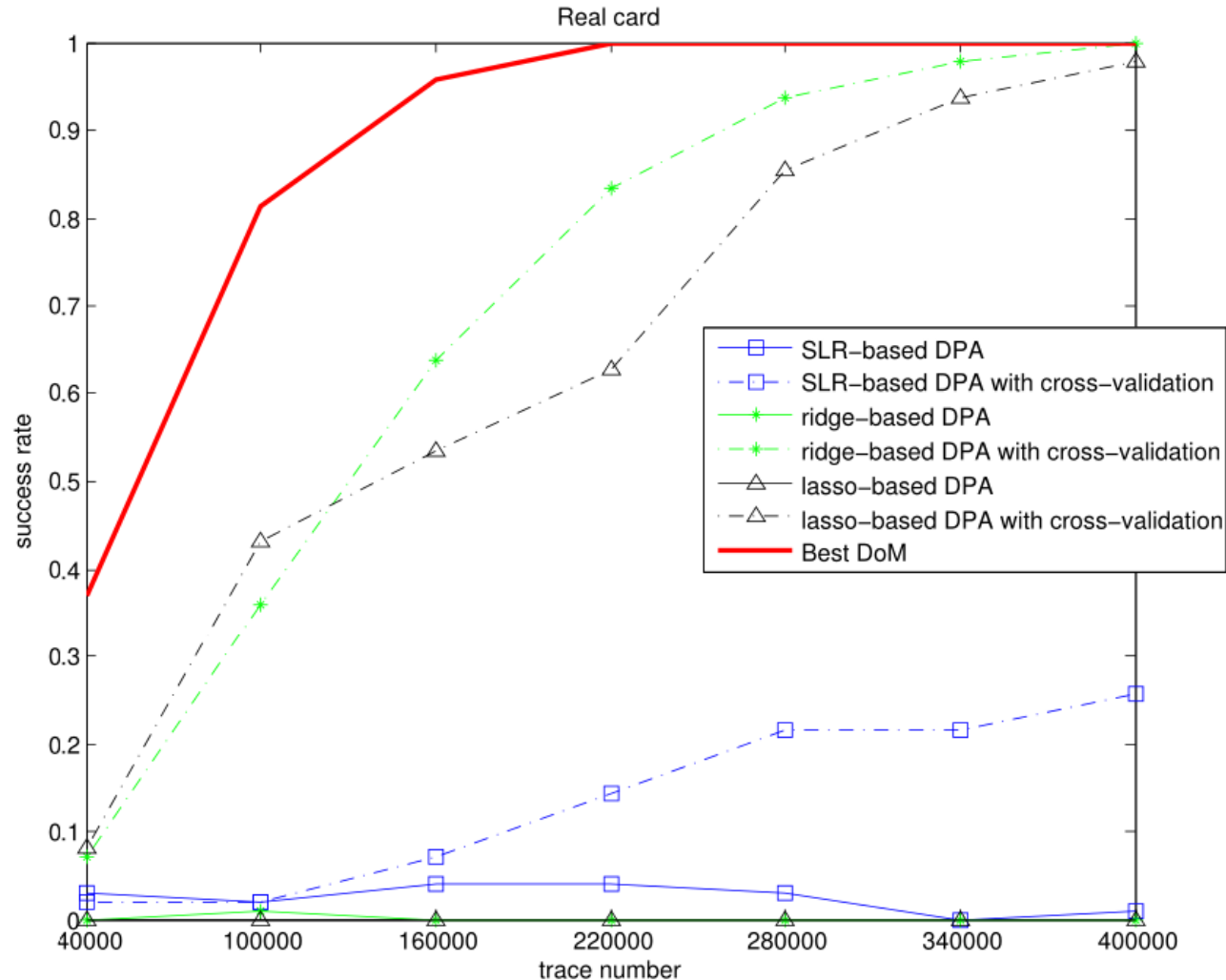
## 4.1.3 Attacks Against Some Artificial Leakage Function

- All low degree terms ( $<4$ ) are discarded.
- Best DoM attack behaves poorly
- The generic-emulating DPAs are not affected.



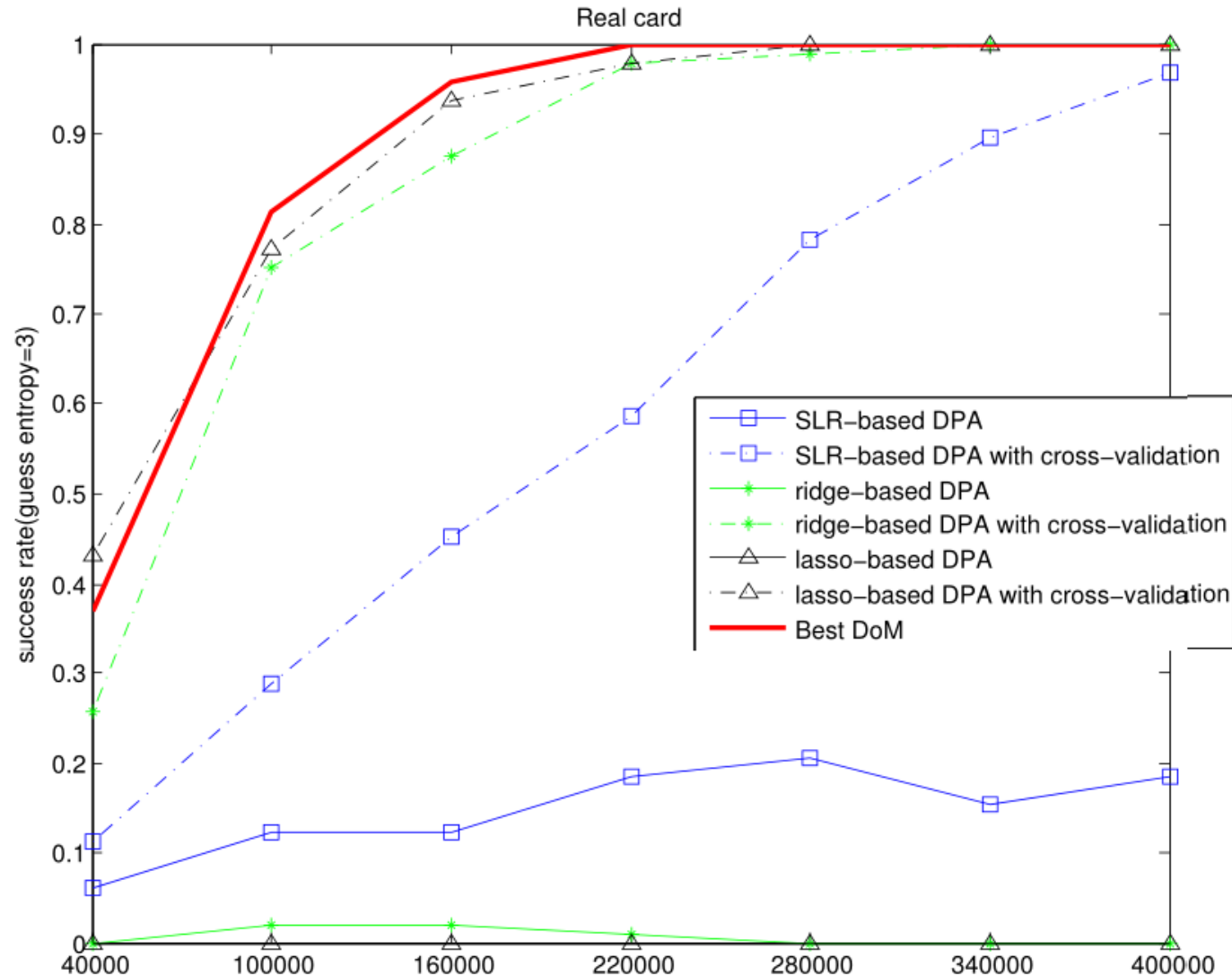
## 4.2 Experiments on Smart Cards

- Microscale ASIC implementation
- 1st order success rates
- C-V significantly improves the performance of generic-emulating DPAs



## 4.2 Experiments on Smart Cards

- 8th-order success rates for better alignment with the best DoM attack
- Ridge-based and lasso-based DPAs (both with C-V) are very close to best DoM.





# Conclusion

## ❖ **Making generic-emulating DPA practicable**

- Ridge-based and lasso-based distinguishers → more stable
- Cross-validation → generic-emulating DPAs can be significantly improved



**Thank you!**